From Classical to Consistent Query Answering under Existential Rules

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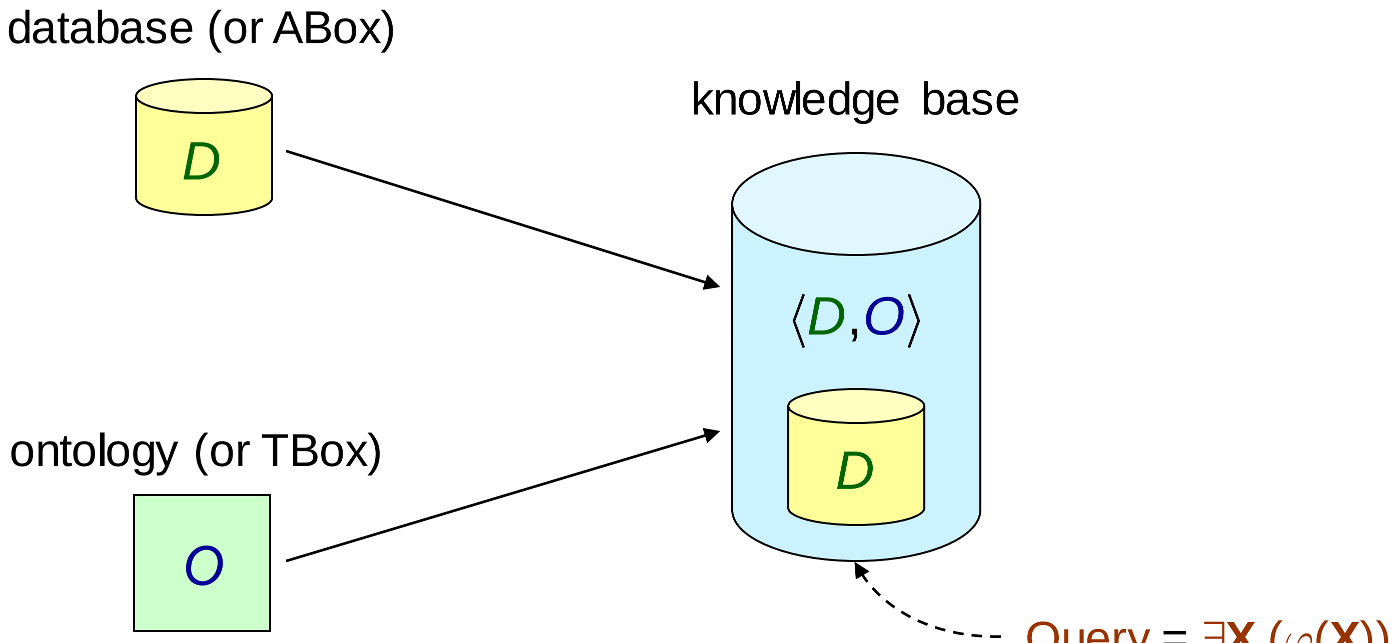
Ontology-based Query Answering (OBQA)



$$\langle D, O\rangle \models \text{Query} \quad \Leftrightarrow \quad D \wedge O \models \text{Query}$$

A Simple Example

D =

professor(John)

fellow(John)

O =

$\forall X\ (professor(X) \rightarrow \exists Y\ (faculty(X) \wedge teaches(X,Y)))$

$\forall X\ (fellow(X) \rightarrow faculty(X))$

$\forall M \models \langle D,O \rangle\ :\ M$ = **...** *teaches*(John,#) **...**

{John $\rightarrow$ John, $X \rightarrow$ #}

$\exists X\ (teaches(\text{John},X))$ ✓

A Simple Example

D =

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$\forall X\ (professor(X) \rightarrow \exists Y\ (faculty(X) \land teaches(X,Y)))$

$\forall X\ (fellow(X) \rightarrow faculty(X))$

$\forall X\ (professor(X) \land fellow(X) \rightarrow \bot)$

no model $\Rightarrow$ every query is entailed

Handling Data Inconsistencies

- The data are likely to be inconsistent with the ontology
- Standard semantics fails: everything is inferred - not meaningful answers
- Two approaches to inconsistency-handling:
  - Resolve the inconsistencies - ideal, but not always possible
  - Live with the inconsistencies - inconsistency-tolerant semantics

ABox Repair (AR) Semantics

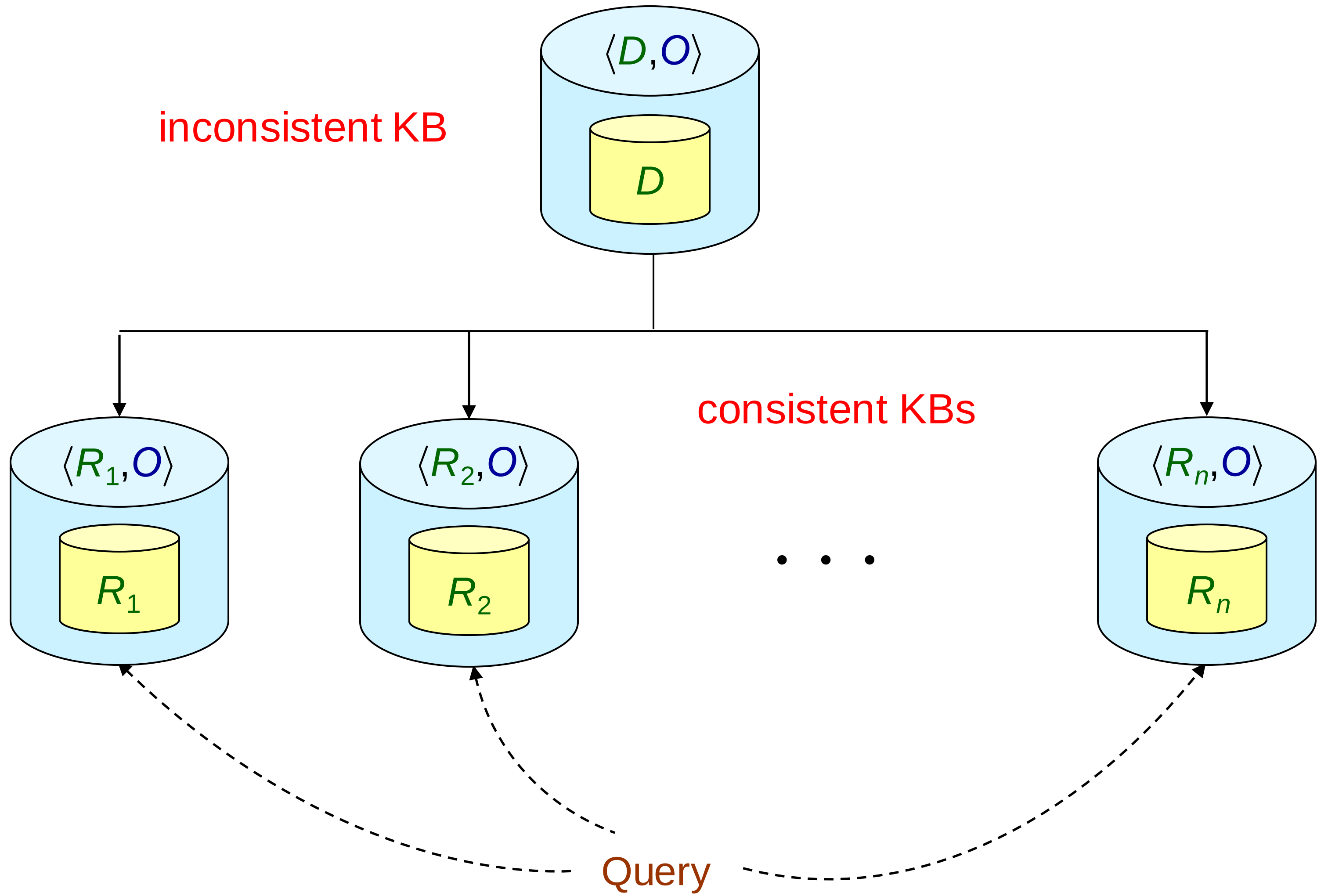
- Standard inconsistency-tolerant semantics

- **IDEA:** The query must be entailed by every database repair

  ⊆-maximal consistent subsets of the database

[Lembo et al., RR 2010]

ABox Repair (AR) Semantics



$\langle D,O\rangle \models_{AR}$ Query $\Leftrightarrow$ $\forall R \in \{R_1,\ldots,R_n\}$: $\langle R,O\rangle \models$ Query

ABox Repair (AR) Semantics: Example

D =

professor(John)

fellow(John)

O =

$\forall X\ (professor(X) \rightarrow \exists Y\ (faculty(X) \wedge teaches(X,Y)))$

$\forall X\ (fellow(X) \rightarrow faculty(X))$

$\forall X\ (professor(X) \wedge fellow(X) \rightarrow \bot)$

$\langle D,O\rangle \models_{\mathrm{AR}} faculty(\mathrm{John})$ ✓

R_1 = *professor*(John)

$\langle R_1,O\rangle \models faculty(\mathrm{John})$ ✓

R_2 = *fellow*(John)

$\langle R_2,O\rangle \models faculty(\mathrm{John})$ ✓

ABox Repair (AR) Semantics: Example

D =

professor(John)

fellow(John)

O =

$\forall X\ (professor(X) \rightarrow \exists Y\ (faculty(X) \wedge teaches(X,Y)))$

$\forall X\ (fellow(X) \rightarrow faculty(X))$

$\forall X\ (professor(X) \wedge fellow(X) \rightarrow \bot)$

$\langle D,O \rangle \models_{\text{AR}} \exists X\ (teaches(\text{John},X))$ ✗

R_1 = *professor*(John)

$\langle R_1,O \rangle \models \exists X\ (teaches(\text{John},X))$ ✓

R_2 = *fellow*(John)

$\langle R_2,O \rangle \models \exists X\ (teaches(\text{John},X))$ ✗

AR Semantics

- Lots of recent work and complexity results for description logics

  [Lembo et al., RR 2010 / Rosati, IJCAI 2011 / Bienvenu, AAAI 2012 / Bienvenu & Rosati, IJCAI 2013]

- This talk is about existential rules + negative constraints

  [Lukasiewicz, Martinez & Simari, ODBASE 2013 / Lukasiewicz, Martinez, P. & Simari, AAAI 2015]

$$\forall \mathbf{X}\ (\varphi(\mathbf{X}) \rightarrow \exists \mathbf{Y}\ (\psi(\mathbf{X},\mathbf{Y}))) \quad + \quad \forall \mathbf{X}\ (\varphi(\mathbf{X}) \rightarrow \bot)$$

Our Goal

Perform an in-depth complexity analysis of consistent query answering under the main classes of existential rules + negative constraints

- Combined
- Bounded-arity combined
- Fixed-program combined
- Data

generic complexity results - from classical to consistent query answering

Combined Complexity

combined or ba-combined or fp-combined   class of ∃-rules   complexity class

$\mathbb{M}$ complexity of classical query answering under $\mathbb{L}$ is $\mathbb{C}$-complete

$$\Downarrow$$

$\mathbb{M}$ complexity of consistent query answering under $\mathbb{L}[\bot]$ is:

$\Pi_{P,2}$-complete   if $\mathbb{C}$ = NP

$\mathbb{C}$-complete   if $\mathbb{C} \supseteq$ PSPACE & $\mathbb{C}$ is deterministic

Combined Complexity: Upper Bounds

Guess and check algorithm (for the complement of the problem)

Input: D, $O \in \mathbb{L}[\bot]$, Q

1. Guess $R \subseteq D$ - a possible repair

2. Verify that R is a repair, i.e., $\langle R,O \rangle$ is consistent and R is $\subseteq$-maximal

3. Verify that $\langle R,O \rangle$ does not entail Q

no harder than classical query answering under $\mathbb{L}$

$\Rightarrow$ our problem is in $\mathrm{coNP}^{\mathbb{C}}$ $\Rightarrow$ in

- $\mathrm{coNP}^{\mathrm{NP}} = \mathrm{co}\Sigma_{P,2} = \Pi_{P,2}$ if $\mathbb{C} = \mathrm{NP}$
- $\mathrm{coNP}^{\mathbb{C}} = \mathrm{co}\mathbb{C} = \mathbb{C}$ if $\mathbb{C} \supseteq \mathrm{PSPACE}$, $\mathbb{C}$ is deterministic

Combined Complexity

combined or ba-combined or fp-combined   class of ∃-rules   complexity class

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A Strong $\Pi_{P,2}$-hardness Result

Consistent query answering under the single constraint

$$\forall X \forall Y \forall Z \forall W \; (p(X,Y,Z) \wedge p(W,X,Z) \rightarrow \bot)$$

while the database and the query use only binary and ternary predicates

(by reduction from satisfiability of 2QBF formulas)

$$\Downarrow$$

For every class $\mathbb{L}$ of existential rules, the fp-combined complexity of consistent query answering under $\mathbb{L}[\bot]$ is $\Pi_{P,2}$-hard

Combined Complexity

combined or ba-combined or fp-combined — class of ∃-rules — complexity class

$\mathbb{M}$ complexity of classical query answering under $\mathbb{L}$ is $\mathbb{C}$-complete

$$\Downarrow$$

$\mathbb{M}$ complexity of consistent query answering under $\mathbb{L}[\bot]$ is:

$\Pi_{P,2}$-complete if $\mathbb{C}$ = NP

$\mathbb{C}$-complete if $\mathbb{C} \supseteq$ PSPACE & $\mathbb{C}$ is deterministic

Data Complexity

class of ∃-rules — complexity class

data complexity of classical query answering under $\mathbb{L}$ is $\mathbb{C}$-complete

$$\Downarrow$$

data complexity of consistent query answering under $\mathbb{L}[\bot]$ is:

coNP-complete if $\mathbb{C} \subseteq$ PTIME

Data Complexity: Upper Bounds

Guess and check algorithm (for the complement of the problem)

Input: D, $O \in \mathbb{L}[\bot]$, Q

1. Guess $R \subseteq D$ - a possible repair
2. Verify that R is a repair, i.e., $\langle R,O \rangle$ is consistent and R is $\subseteq$-maximal
3. Verify that $\langle R,O \rangle$ does not entail Q

no harder than classical query answering under $\mathbb{L}$

$\Rightarrow$ our problem is in $\text{coNP}^{\mathbb{C}}$ $\Rightarrow$ in coNP (since $\text{NP}^{\text{PTIME}} = \text{NP}$)

A Strong coNP-hardness Result

Consistent query answering under the single constraint

$$\forall X \ (p(X) \wedge s(X) \ \rightarrow \ \bot)$$

while the query is fixed

(by reduction from 2+2UNSAT)

$$\Downarrow$$

For every class $\mathbb{L}$ of existential rules, the data complexity of consistent query answering under $\mathbb{L}[\bot]$ is coNP-hard

Data Complexity

class of ∃-rules　　　complexity class

data complexity of classical query answering under $\mathbb{L}$ is $\mathbb{C}$-complete

$$\Downarrow$$

data complexity of consistent query answering under $\mathbb{L}[\bot]$ is:

coNP-complete　　if　　$\mathbb{C} \subseteq$ PTIME

From Classical to Consistent Query Answering

(ba-/fp)combined complexity:

in NP $\rightarrow$ $\Pi_{P,2}$-complete

$\mathbb{C}$-complete, $\mathbb{C} \supseteq$ PSPACE & $\mathbb{C}$ is deterministic $\rightarrow$ $\mathbb{C}$-complete

data complexity:

in $\mathbb{C} \subseteq$ PTIME $\rightarrow$ coNP-complete

an (almost) complete picture for the main classes of existential rules + negative constraints

Existential Rules

$$\forall \mathbf{X}\ (\varphi(\mathbf{X}) \rightarrow \exists \mathbf{Y}\ (\psi(\mathbf{X},\mathbf{Y})))$$

conjunctions of atoms

- Classical query answering under existential rules is undecidable
  see, e.g., [Beeri & Vardi, ICALP 1981]

- Expressive decidable fragments - field of intense research
  (e.g., Montpellier, Dresden, Calabria, Oxford, Vienna, …)

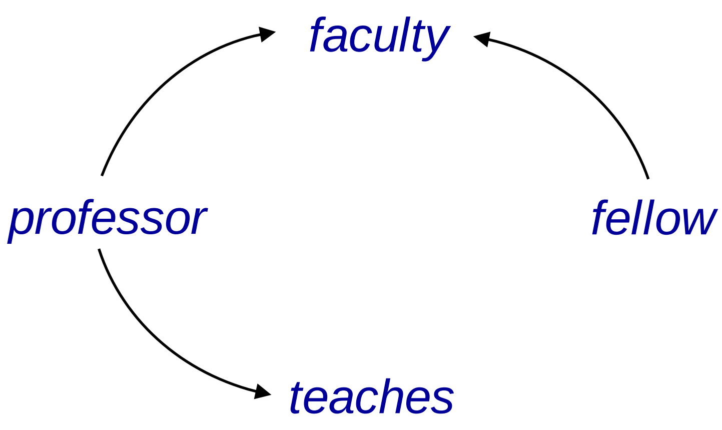
- Main decidability paradigms: acyclicity, guardedness & stickiness

Acyclic Existential Rules

- The predicate graph is acyclic

$$\forall X\ (professor(X) \ \rightarrow\ \exists Y\ (faculty(X) \wedge teaches(X,Y)))$$

$$\forall X\ (fellow(X) \ \rightarrow\ faculty(X))$$



(Frontier-)Guarded Existential Rules

- Frontier-guardedness: There exists a body-atom that contains the frontier

$$\forall X \forall Y \forall Z\ (\boldsymbol{supervisorOf(X,Y)} \wedge supervisorOf(Y,Z) \rightarrow manager(X))$$

- Guardedness: There exists a body-atom that contains all the ∀-variables

$$\forall X \forall Y\ (\boldsymbol{supervisorOf(X,Y)} \wedge emp(Y) \rightarrow emp(X))$$

- Linearity: There exists only one atom in the body

$$\forall X\ (\boldsymbol{employee(X)} \rightarrow \exists Y\ (supervisorOf(Y,X) \wedge employee(Y)))$$

Sticky Existential Rules

- Join-variables stick to the inferred atoms

$$\forall X \forall Y \forall Z\, (q(X,Y) \wedge p(Y,Z) \rightarrow \exists W\, (t(X,Y,W)))$$

$$\forall X \forall Y \forall Z\, (t(X,Y,Z) \rightarrow \exists W\, (s(Y,W)))$$

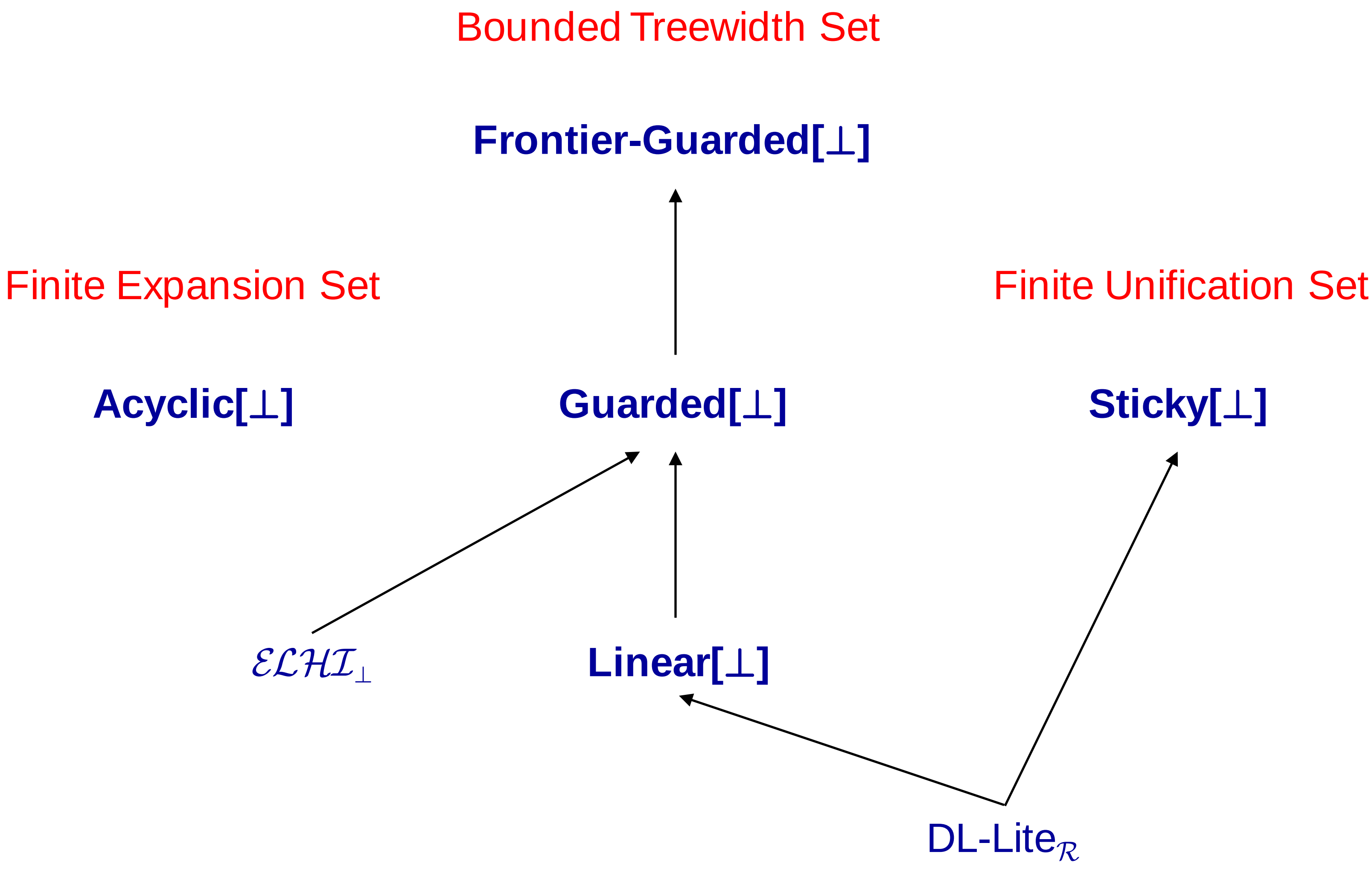
✓

$$\forall X \forall Y \forall Z\, (q(X,Y) \wedge p(Y,Z) \rightarrow \exists W\, (t(X,Y,W)))$$

$$\forall X \forall Y \forall Z\, (t(X,Y,Z) \rightarrow \exists W\, (s(X,W)))$$

✗

Existential Rules + Negative Constraints



From Classical to Consistent Query Answering

(ba-/fp)combined complexity:

in NP $\rightarrow$ $\Pi_{p,2}$-complete

$\mathbb{C}$-complete, $\mathbb{C} \supseteq$ PSPACE & $\mathbb{C}$ is deterministic $\rightarrow$ $\mathbb{C}$-complete

data complexity:

in $\mathbb{C} \subseteq$ PTIME $\rightarrow$ coNP-complete

we simply need to exploit existing results on classical query answering

Classical Query Answering

	Combined	ba-combined	fp-combined	Data
Acyclic[⊥]	NEXPTIME	NEXPTIME	NP	in AC_0
Frontier-Guarded[⊥]	2EXPTIME	2EXPTIME	NP	PTIME
Guarded[⊥]	2EXPTIME	EXPTIME	NP	PTIME
Linear[⊥]	PSPACE	NP	NP	in AC_0
Sticky[⊥]	EXPTIME	NP	NP	in AC_0

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Frontier-Guarded[⊥]	2EXPTIME	2EXPTIME	NP	PTIME
Guarded[⊥]	2EXPTIME	EXPTIME	NP	PTIME
Linear[⊥]	PSPACE	NP	NP	in AC_0
Sticky[⊥]	EXPTIME	NP	NP	in AC_0

- Until recently, it was generally believed that it is EXPTIME
- The obvious algorithm does not work - models of double-exponential size

Classical Query Answering

	Combined	ba-combined	fp-combined	Data
Acyclic[⊥]	NEXPTIME	NEXPTIME	NP	in AC_0
Frontier-Guarded[⊥]	2EXPTIME	2EXPTIME	NP	PTIME
Guarded[⊥]	2EXPTIME	EXPTIME	NP	PTIME
Linear[⊥]	PSPACE	NP	NP	in AC_0
Sticky[⊥]	EXPTIME	NP	NP	in AC_0

- Upper bound: non-deterministically construct a proof of the query
- Lower bound: by reduction from a TILING problem

Classical Query Answering

	Combined	ba-combined	fp-combined	Data
Acyclic[⊥]	NEXPTIME	NEXPTIME	NP	in AC_0
Frontier-Guarded[⊥]	2EXPTIME	2EXPTIME	NP	PTIME
Guarded[⊥]	2EXPTIME	EXPTIME	NP	PTIME
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Sticky[⊥]	EXPTIME	NP	NP	in AC_0

(ba-/fp)combined complexity:

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$\mathbb{C}$-complete, $\mathbb{C} \supseteq$ PSPACE & $\mathbb{C}$ is deterministic $\rightarrow$ $\mathbb{C}$-complete

data complexity:

in $\mathbb{C} \subseteq$ PTIME $\rightarrow$ coNP-complete

Consistent Query Answering

	Combined	ba-combined	fp-combined	Data
Acyclic[⊥]	?	?	$\Pi_{P,2}$	coNP
Frontier-Guarded[⊥]	2EXPTIME	2EXPTIME	$\Pi_{P,2}$	coNP
Guarded[⊥]	2EXPTIME	EXPTIME	$\Pi_{P,2}$	coNP
Linear[⊥]	PSPACE	$\Pi_{P,2}$	$\Pi_{P,2}$	coNP
Sticky[⊥]	EXPTIME	$\Pi_{P,2}$	$\Pi_{P,2}$	coNP

(ba-/fp)combined complexity:

in NP $\rightarrow$ $\Pi_{P,2}$-complete

$\mathbb{C}$-complete, $\mathbb{C} \supseteq$ PSPACE & $\mathbb{C}$ is deterministic $\rightarrow$ $\mathbb{C}$-complete

data complexity:

in $\mathbb{C} \subseteq$ PTIME $\rightarrow$ coNP-complete

Complexity of Acyclic[$\perp$]

- The guess and check algorithm gives a coNP$^{\text{NEXPTIME}}$ upper bound
- The class NP$^{\text{NEXPTIME}}$ lies at a higher level of the strong exponential hierarchy
- The SEH collapses to its Δ_2 level $\Rightarrow$ NP$^{\text{NEXPTIME}}$ = P$^{\text{NE}}$
  [Hemachandra, J. Comput. Syst. Sci. 1989]
- P$^{\text{NE}}$ is a deterministic class $\Rightarrow$ coP$^{\text{NE}}$ = P$^{\text{NE}}$

Consistent Query Answering

	Combined	ba-combined	fp-combined	Data
Acyclic[⊥]	NEXP - P^{NE}	NEXP - P^{NE}	$\Pi_{P,2}$	coNP
Frontier-Guarded[⊥]	2EXPTIME	2EXPTIME	$\Pi_{P,2}$	coNP
Guarded[⊥]	2EXPTIME	EXPTIME	$\Pi_{P,2}$	coNP
Linear[⊥]	PSPACE	$\Pi_{P,2}$	$\Pi_{P,2}$	coNP
Sticky[⊥]	EXPTIME	$\Pi_{P,2}$	$\Pi_{P,2}$	coNP

$P^{NE} \subseteq coNEXPTIME^{NP}$

[Hemachandra, J. Comput. Syst. Sci. 1989]

Consistent Query Answering

	Combined	ba-combined	fp-combined	Data
Acyclic[⊥]	NEXP - P^{NE}	NEXP - P^{NE}	$\Pi_{P,2}$	coNP
Frontier-Guarded[⊥]	2EXPTIME	2EXPTIME	$\Pi_{P,2}$	coNP
Guarded[⊥]	2EXPTIME	EXPTIME	$\Pi_{P,2}$	coNP
Linear[⊥]	PSPACE	$\Pi_{P,2}$	$\Pi_{P,2}$	coNP
Sticky[⊥]	EXPTIME	$\Pi_{P,2}$	$\Pi_{P,2}$	coNP

Conjecture: Consistent query answering under Acyclic[⊥] is coNEXPTIME^{NP}-c

Data Intractable

	Combined	ba-combined	fp-combined	Data
Acyclic[⊥]	NEXP - P^{NE}	NEXP - P^{NE}	$\Pi_{P,2}$	coNP
Frontier-Guarded[⊥]	2EXPTIME	2EXPTIME	$\Pi_{P,2}$	coNP
Guarded[⊥]	2EXPTIME	EXPTIME	$\Pi_{P,2}$	coNP
Linear[⊥]	PSPACE	$\Pi_{P,2}$	$\Pi_{P,2}$	coNP
Sticky[⊥]	EXPTIME	$\Pi_{P,2}$	$\Pi_{P,2}$	coNP

but, what about tractability results w.r.t. the data complexity?

…consider approximations of the AR semantics

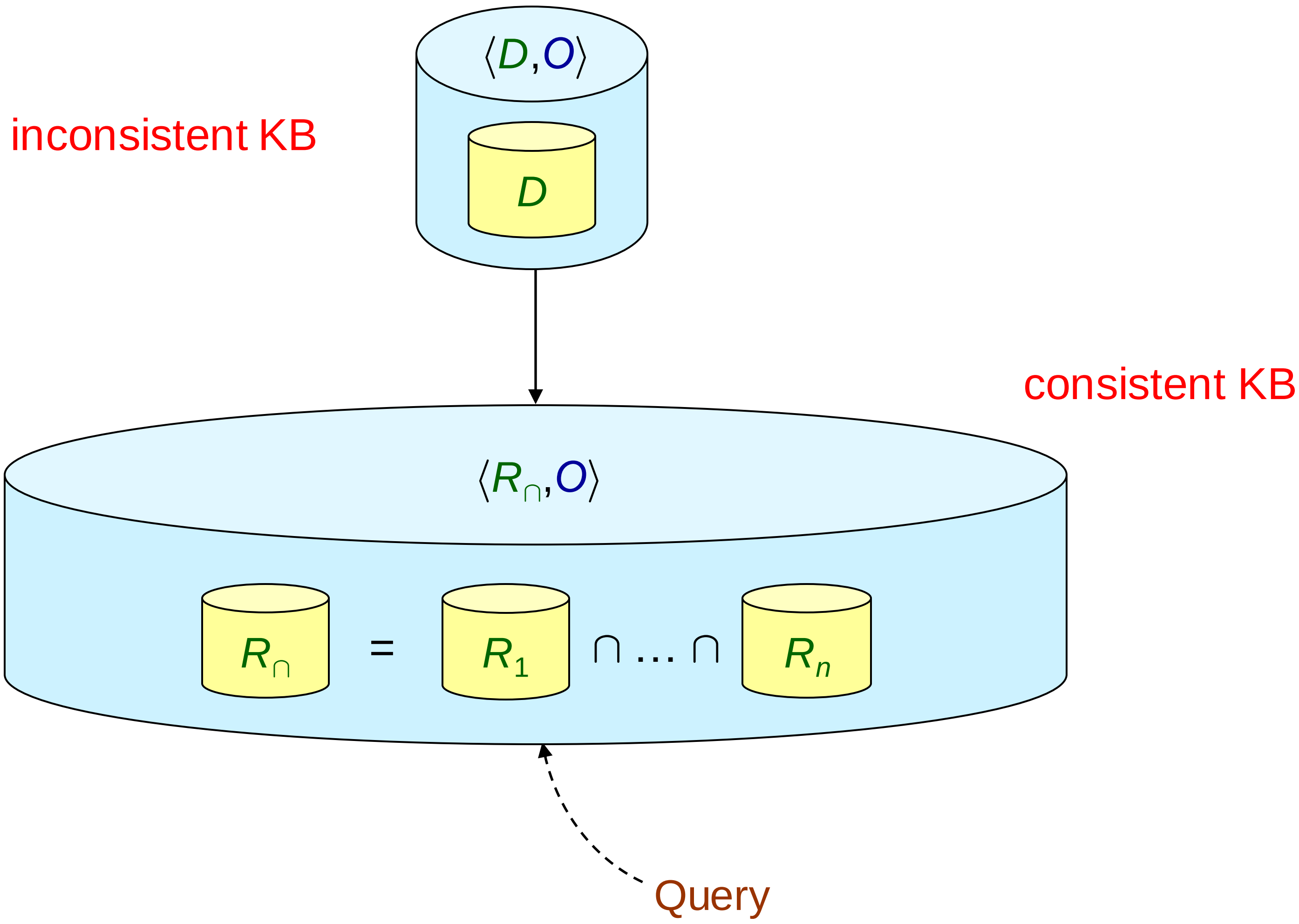
Intersection ABox Repair (IAR) Semantics

- One of the basic sound approximations of the AR semantics
- **IDEA:** The query must be entailed by the intersection of the database repairs

↙ $\subseteq$-maximal consistent subsets of the database

[Lembo et al., RR 2010]

Intersection ABox Repair (IAR) Semantics



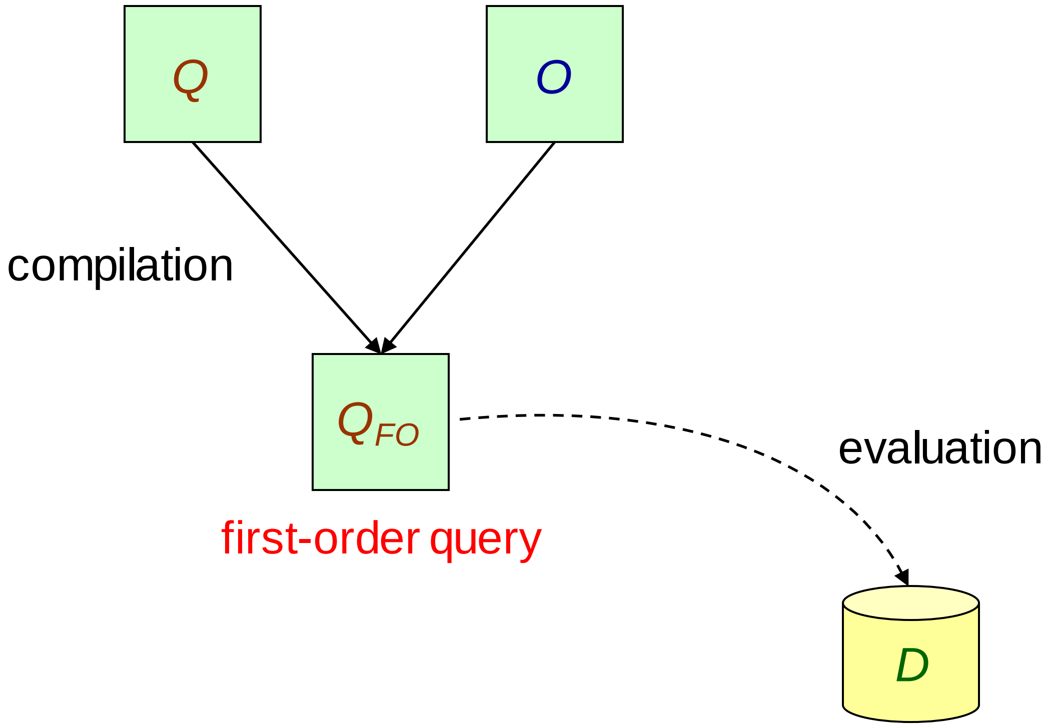
$\langle D, O \rangle \models_{IAR}$ Query $\Leftrightarrow$ $\langle R_\cap, O \rangle \models$ Query

Data Complexity under the IAR Semantics

Acyclic[⊥]	in AC_0
Frontier-Guarded[⊥]	coNP
Guarded[⊥]	coNP
Linear[⊥]	in AC_0
Sticky[⊥]	in AC_0

via first-order rewritability - a generic result can be established

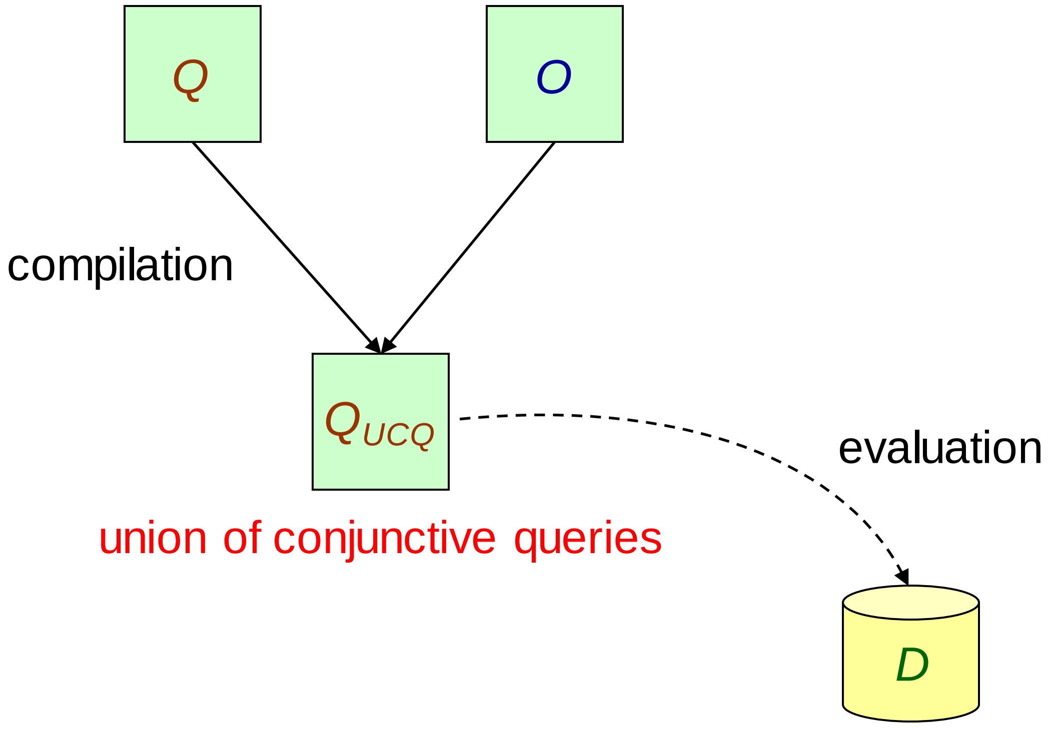
First-Order Rewritability (FO-Rewritability)



$$\forall D \; : \; \langle D, O \rangle \models Q \quad \Leftrightarrow \quad D \models Q_{FO}$$

$$\forall D \; : \; \langle D, O \rangle \models_{IAR} Q \quad \Leftrightarrow \quad D \models Q_{FO}$$

UCQ-Rewritability



$$\forall D \;:\; \langle D, O \rangle \models Q \quad \Leftrightarrow \quad D \models Q_{UCQ}$$

$$\forall D \;:\; \langle D, O \rangle \models_{\mathrm{IAR}} Q \quad \Leftrightarrow \quad D \models Q_{UCQ}$$

From UCQ-Rewritability to FO-Rewritability

class of ∃-rules

classical query answering under $\mathbb{L}$ is UCQ-Rewritable

$$\Downarrow$$

consistent query answering under the IAR semantics for $\mathbb{L}[\perp]$ is FO-Rewritable

Data Complexity under the IAR Semantics

Acyclic[⊥]	in AC_0
Frontier-Guarded[⊥]	coNP
Guarded[⊥]	coNP
Linear[⊥]	in AC_0
Sticky[⊥]	in AC_0

via first-order rewritability - a generic result can be established

Key Message

We can transfer complexity results

from classical to consistent query answering

in a generic and uniform way

…with some unexpected exceptions - Acyclic[$\perp$]

Thank you!